

Number Representation and Binary Arithmetic

Unsigned Addition

Binary Addition

$$0 + 0 = 0$$

$$0 + 1 = 1$$

$$1 + 0 = 1$$

$$1 + 1 = 0 \text{ and carry } 1 \text{ to the next column}$$

Examples:

$$\begin{array}{r} 0101 \\ + 0010 \\ \hline 0111 \end{array} \quad \begin{array}{l} (5_{10}) \\ (2_{10}) \\ (7_{10}) \end{array}$$

$$\begin{array}{r} \overset{1 \ 1 \ 1}{0101} \\ + \ 0011 \\ \hline 1000 \end{array} \quad \begin{array}{l} (5_{10}) \\ (3_{10}) \\ (8_{10}) \end{array} \quad \begin{array}{c} \longleftarrow \text{Carries} \end{array}$$

Binary Addition w/Overflow

Add 45_{10} and 44_{10} in binary

$$\begin{array}{r} \begin{array}{ccccccc} & & 1 & & 1 & 1 & \\ & & \leftarrow & & & & \end{array} & \text{Carries} \\ 101101 & (45_{10}) \\ + 101100 & (44_{10}) \\ \hline 1011001 & (89_{10}) \end{array}$$

If the operands are unsigned, you can use the final carry-out as the MSB of the result.

Adding 2 k bit numbers $\Leftrightarrow k+1$ bit result

More Binary Addition w/Overflow

$$\begin{array}{r} \overset{1}{1} \overset{1}{1} \overset{1}{1} \overset{1}{1} \\ 1111 \\ + 0001 \\ \hline 0000 \end{array} \quad \begin{array}{l} (15_{10}) \\ (1_{10}) \\ (0_{10}) \end{array}$$

If you don't want a 5-bit result, just keep the lower 4 bits.

This is insufficient to hold the result (16).

It rolls over back to 0.

Signed Numbers

Negative Binary Numbers

- Several ways of representing negative numbers
- Most obvious is to add a sign (+ or -) to the binary integer

Sign-Magnitude

0 is +

1 is -

Number	Sign	Magnitude	Full Number
+1	0	0001	00001
-1	1	0001	10001
+5	0	0101	00101
-5	1	0101	10101
+0	0	0000	00000
-0	1	0000	10000

Easy to interpret number value

Sign-Magnitude Examples

$$\begin{array}{r}
 \begin{array}{ccc} & 0 & 0 & 0 \\ 0 & 0011 & (+3_{10}) \\ + & 0 & 0100 & (+4_{10}) \\ \hline 0 & 0111 & (+7_{10}) \end{array}
 \end{array}$$

$$\begin{array}{r}
 \begin{array}{ccc} & 0 & 0 & 0 \\ 1 & 0011 & (-3_{10}) \\ + & 1 & 0100 & (-4_{10}) \\ \hline 1 & 0111 & (-7_{10}) \end{array}
 \end{array}$$

Signs are the same

Just add the magnitudes

Another Sign-Magnitude Example

Signs are different -
determine which is
larger magnitude

$$\begin{array}{rcl}
 & 0 & 0101 & (+5_{10}) \\
 + & 1 & 0011 & (-3_{10}) \\
 \hline
 \end{array}$$

Put larger magnitude
number on top

$$\begin{array}{rcl}
 & 0 & 0101 & (+5_{10}) \\
 - & 1 & 0011 & (-3_{10}) \\
 \hline
 & 0 & 0010 & (+2_{10})
 \end{array}$$

Subtract

Result has sign of larger
magnitude number...

Yet Another Sign-Magnitude Example

Signs are different -
determine which is
larger magnitude

$$\begin{array}{rcl}
 & 0 & 0010 & (+2_{10}) \\
 + & 1 & 0101 & (-5_{10}) \\
 \hline
 \end{array}$$

Put larger magnitude
number on top

$$\begin{array}{rcl}
 & 1 & 0101 & (-5_{10}) \\
 - & 0 & 0010 & (+2_{10}) \\
 \hline
 & 1 & 0011 & (-3_{10})
 \end{array}$$

Subtract

Result has sign of larger
magnitude number...

Sign-Magnitude

- Addition requires two separate operations
 - addition
 - subtraction
- Several decisions:
 - Signs same or different?
 - Which operand is larger?
 - What is sign of final result?
- Two zeroes (+0, -0)

Sign-Magnitude

- Advantages
 - Easy to understand
- Disadvantages
 - Two different 0s
 - Hard to implement in logic

One's Complement

- Positive numbers are the same as sign-magnitude
- $-N$ is represented as the complement of N
+ $-N = N'$

Number	1's Complement
+1	0001
-1	1110
+5	0101
-5	1010
+0	0000
-0	1111

One's Complement Examples

$$\begin{array}{r}
 000101 \\
 + 010100 \\
 \hline
 011001
 \end{array}
 \begin{array}{l}
 (5_{10}) \\
 (20_{10}) \\
 (25_{10})
 \end{array}$$

$$\begin{array}{r}
 \boxed{1100} \\
 0101 \\
 + 1100 \\
 \hline
 0001 \leftarrow + \\
 \hline
 0010
 \end{array}
 \begin{array}{l}
 (5_{10}) \\
 (-3_{10}) \\
 (+2_{10})
 \end{array}$$

If there is a carry out on the left, it must be *wrapped around* and added back in on the right

One's Complement

- Addition complicated by end around carry
- No decisions (like SM)
- Still two zeroes (+0, -0)

One's Complement

- Advantages
 - Easy to generate $-N$
 - Only one addition process
- Disadvantages
 - End around carry
 - Two different 0s

Two's Complement

- Treat positional digits differently

$$0111_{2C} = -0 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 = 7_{10}$$

$$1111_{2C} = -1 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 = -1_{10}$$


Most significant bit (MSB) given negative weight...
Other bits same as in unsigned

Two's Complement

Number	2's Complement
+1	0001
-1	1111
+5	0101
-5	1011
+0	0000
-0	none

Sign-Extension in 2's Complement

- To make a k -bit number wider
 - replicate sign bit

$$110_{2C} = -1 \times 2^2 + 1 \times 2^1 = -2_{10}$$

$$1110_{2C} = -1 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 = -2_{10}$$

$$11110_{2C} = -1 \times 2^4 + 1 \times 2^3 + 1 \times 2^2 + 1 \times 2^1 = -2_{10}$$

More Sign-Extension

$$010_{2C} = 1 \times 2^1 = 2_{10}$$

$$0010_{2C} = 1 \times 2^1 = 2_{10}$$

$$00010_{2C} = 1 \times 2^1 = 2_{10}$$

Works for both positive and negative numbers

Negating a 2's Complement Number

- Invert all the bits
- Add 1

invert

$$+2_{10} = 0010_{2C} \rightarrow 1101 + 1 = 1110_{2C} = -2_{10}$$

invert

$$-2_{10} = 1110_{2C} \rightarrow 0001 + 1 = 0010_{2C} = +2_{10}$$

Two's Complement Addition

$$\begin{array}{r} 01 \\ 0101 \quad (+5_{10}) \\ + 0001 \quad (+1_{10}) \\ \hline 0110 \quad (+6_{10}) \end{array}$$

$$\begin{array}{r} 11 \\ 1011 \quad (-5_{10}) \\ + 1111 \quad (-1_{10}) \\ \hline 1010 \quad (-6_{10}) \end{array}$$

$$\begin{array}{r} 11 \\ 0110 \quad (+6_{10}) \\ + 1111 \quad (-1_{10}) \\ \hline 0101 \quad (+5_{10}) \end{array}$$

Operation is same as for unsigned
Same rules, same procedure

Interpretation of operands and
results are different

Two's Complement

- Addition always the same
- Only 1 zero
- Representation of choice...

Two's Complement Overflow

- All representations can overflow
 - Focus on 2's complement here

$$\begin{array}{rcl} & 1 & 0 & 0 & 0 \\ & 1 & 0 & 1 & 1 & (-5_{10}) \\ + & 1 & 1 & 0 & 0 & (-4_{10}) \\ \hline & 0 & 1 & 1 & 1 & (+7_{10}) \quad ?? \end{array}$$

The correct answer (-9) cannot be represented by 4 bits
The correct answer is 10111

Overflow

- Can you use leftmost carry-out as new MSB?

$$\begin{array}{r} 1\ 0\ 0\ 0 \\ 1011 \\ + 1100 \\ \hline 10111 \end{array} \quad \begin{array}{l} (-5_{10}) \\ (-4_{10}) \\ (-9_{10}) \end{array}$$

Works here...

$$\begin{array}{r} 0\ 0\ 0\ 0 \\ 1000 \\ + 0111 \\ \hline 01111 \end{array} \quad \begin{array}{l} (-8_{10}) \\ (+7_{10}) \\ (+15_{10}) \end{array} \quad ??$$

Does NOT work here...

- The answer is no, in general

Handling Overflow

1. Sign-extend the operands
2. Do the addition

$$\begin{array}{rcl} & 0 & 0 & 0 & 0 \\ & \textcolor{red}{1} & 1 & 0 & 0 & 0 & (-8_{10}) \\ + & \textcolor{red}{0} & 0 & 1 & 1 & 1 & (+7_{10}) \\ \hline & 1 & 1 & 1 & 1 & 1 & (-1_{10}) \end{array}$$

When Can Overflow Occur?

- Adding two positive numbers - yes
- Adding two negative numbers - yes
- Adding a positive to a negative - no
- Adding a negative to a positive - no

General Handling of Overflow

- Adding two k bit numbers gives $k+1$ bit result

Two options include:

1. Sign extend and keep entire $k+1$ bit result
 2. Throw away new bit generated (keep only k bits)
 - i. Detect overflow and signal an error
 - ii. Ignore the overflow (this is what most computers do)
- The choice is up to the designer

Detecting Overflow

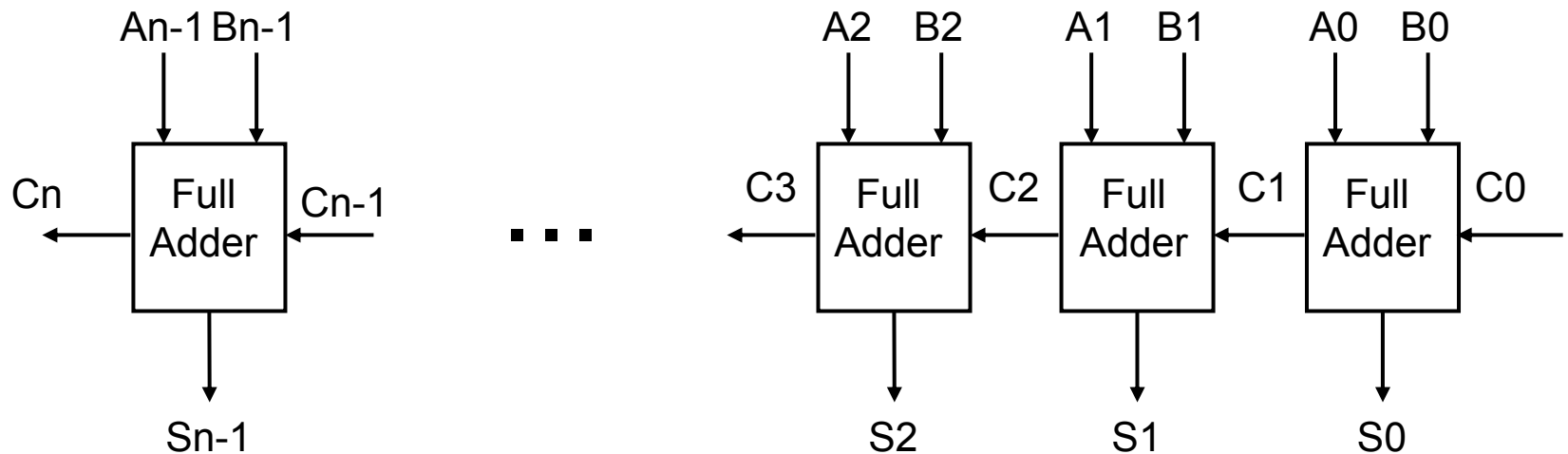
- Overflow occurs if:
 - Adding two positive numbers gives what looks like a negative result
 - Adding two negative numbers gives what looks like a positive result
- Overflow will never occur when adding positive to negative

There are other ways of detecting overflow vs. what is suggested here...

Arithmetic Circuits

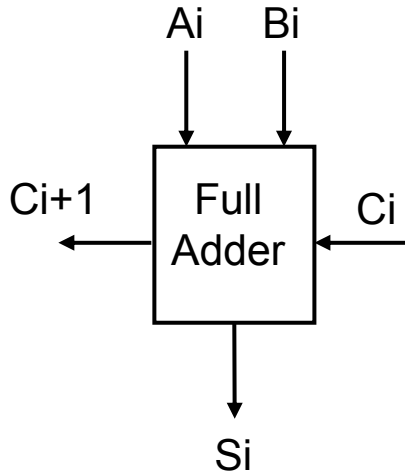
Binary Adder

- An example of an iterative network

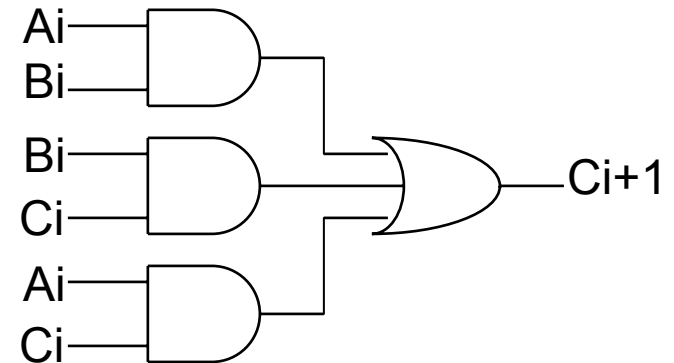
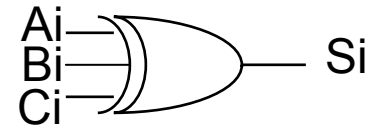


This type of adder is also called a ripple-carry adder because the carry ripples through from cell to cell

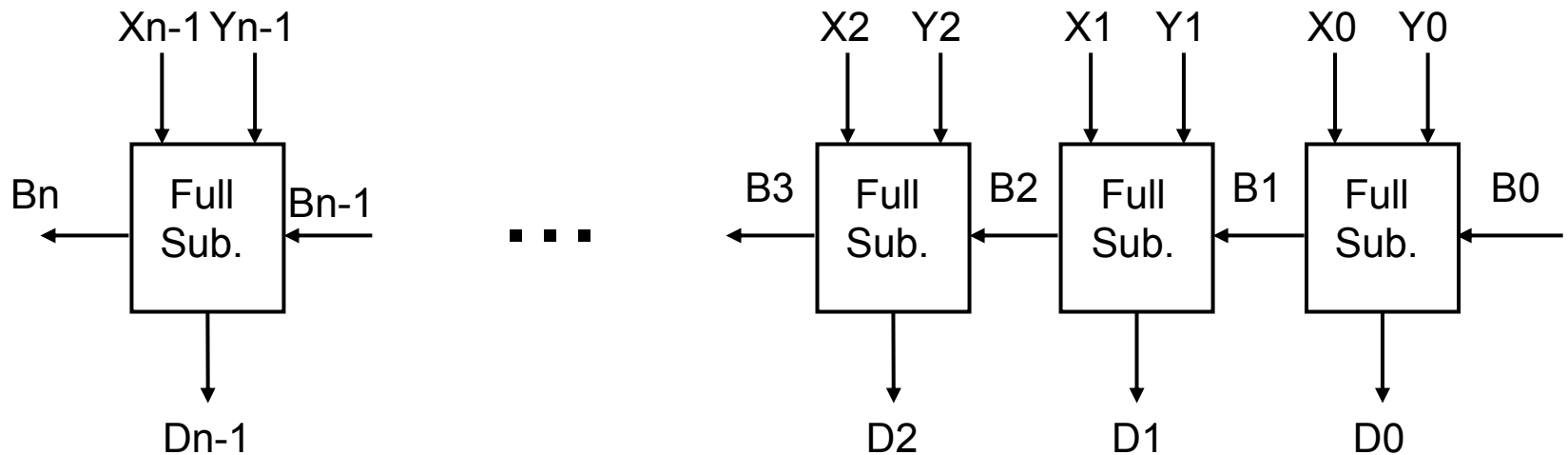
Full Adder Derivation



Ai	Bi	Ci	Ci+1	Si
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1



Binary Subtractor

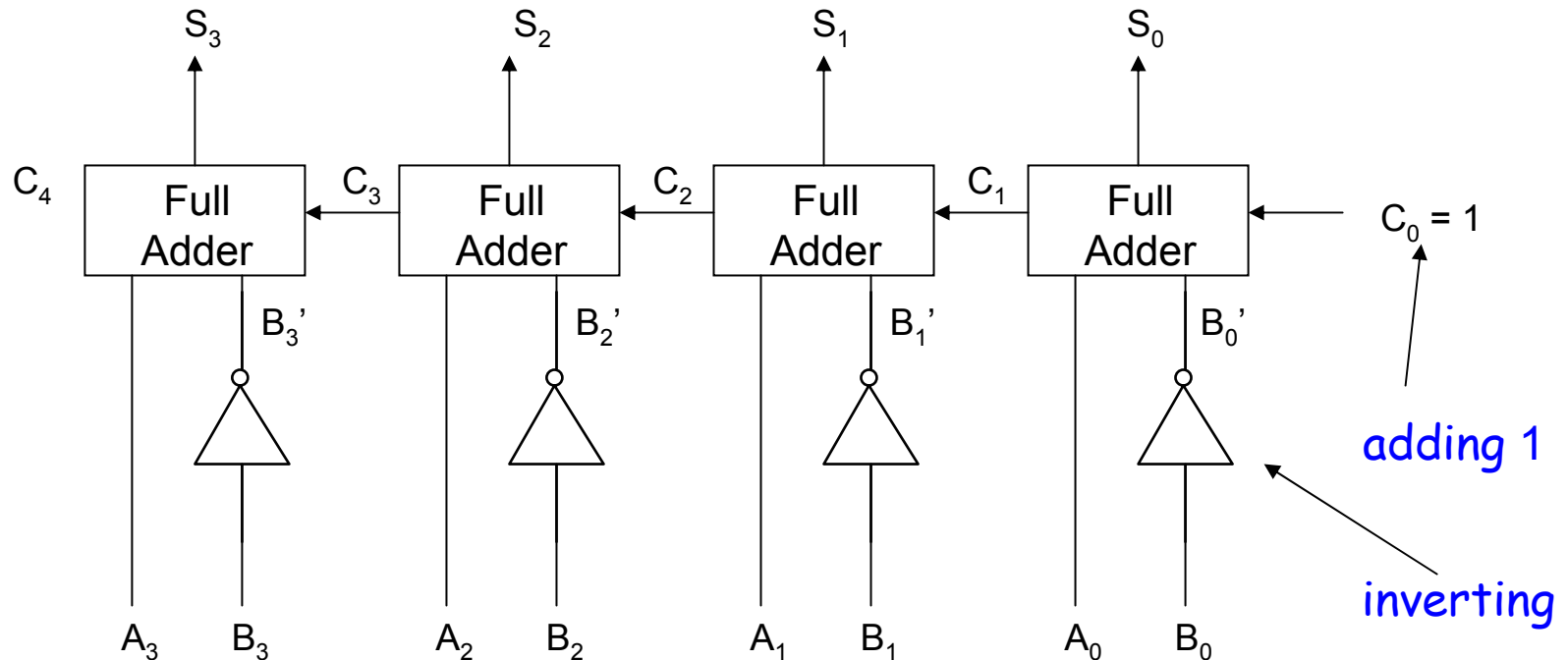


$D(\text{ifference}) = X - Y$

B is the borrow signal

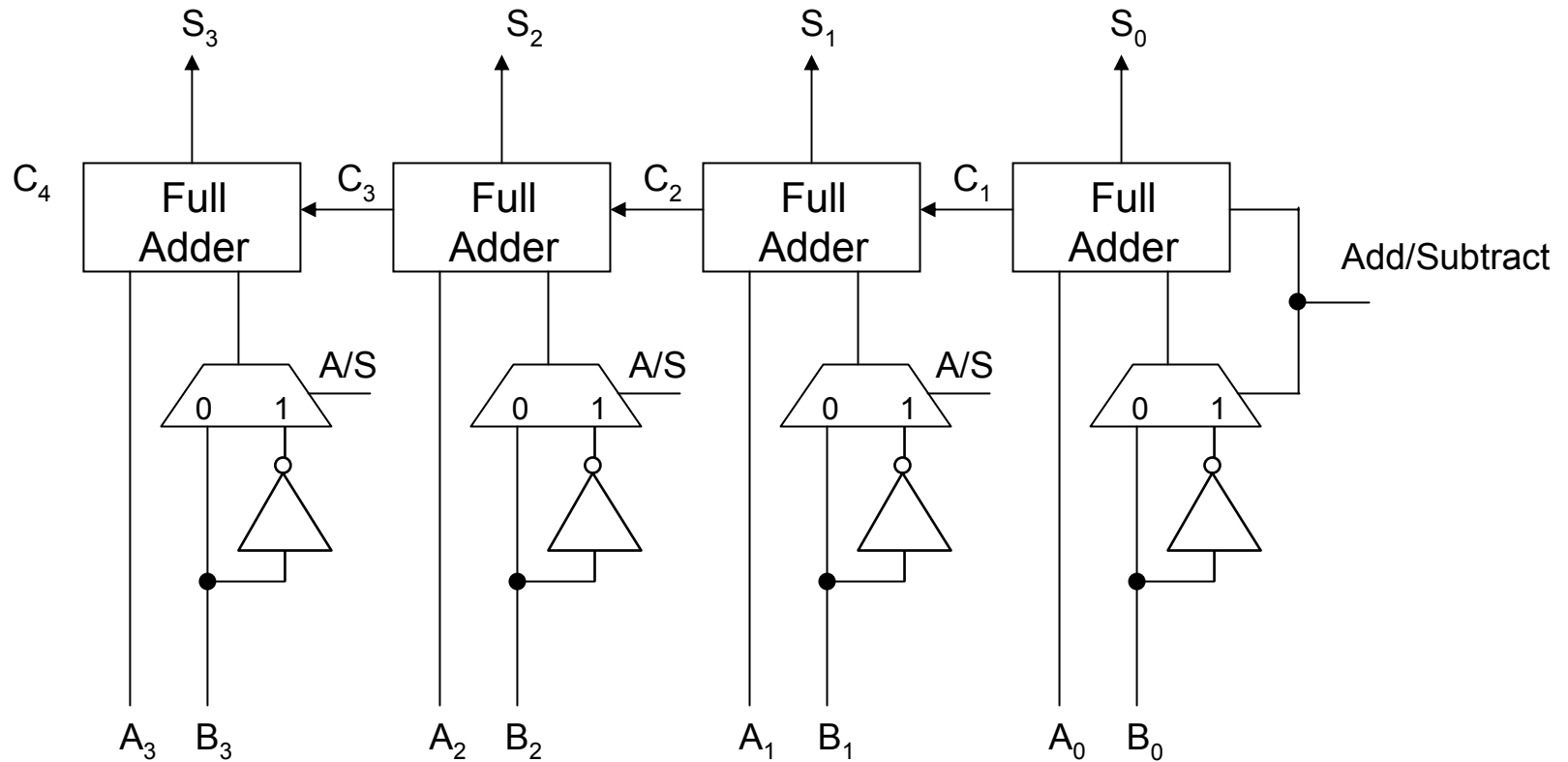
Equations generated
similarly to full adder

Another Way to Make a Subtractor (2's complement)



$S = A + (-B)$
-B is generated by inverting and adding 1

Binary Adder/Subtractor (2's complement)



Binary Arithmetic Comparison

	Sign Magnitude	One's Complement	Two's Complement
Negative Number	Easiest to Understand Simple to Compute	Easy to Compute	Hardest to Compute
Zeroes	2 Zeroes	2 Zeroes	1 Zero
Largest Number	Same number of + and - Numbers	Same number of + and - Numbers	One Extra Negative Number
Logic Required	Requires Adder and Subtractor	Only Adder Required	Only Adder Required
	Extra Logic to Identify Larger Operand, Compute Sign, etc.	Carry Wraps Around	-
Overflow Detection	Overflow: Carry from High Order Adder Bits	Overflow: Sign of Both Operands is the Same and Sign of Sum is Different	Overflow: Sign of Both Operands is the Same and Sign of Sum is Different

Summary

- Understand the three main representation of binary integers
 - Sign-Magnitude
 - One's Complement
 - Two's Complement
- Design binary adders and subtracters for each representation