

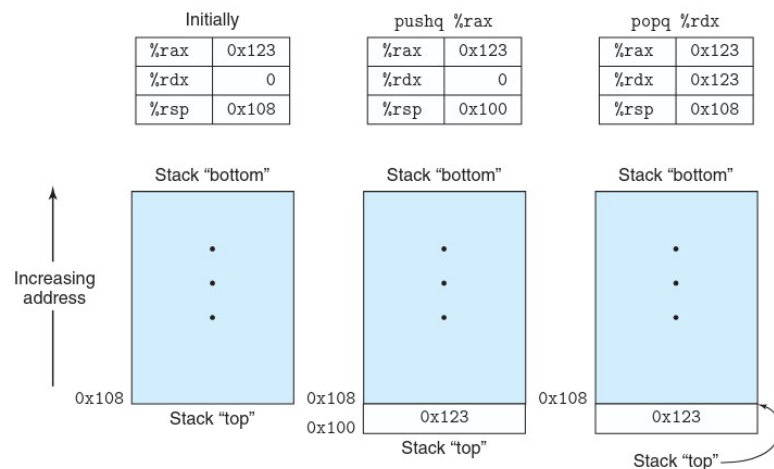
Run-Time Stack (ch 3.4.4)

- pushX
- popX

Instruction	Effect	Description
pushq <i>S</i>	$R[\%rsp] \leftarrow R[\%rsp] - 8;$ $M[R[\%rsp]] \leftarrow S$	Push quad word
popq <i>D</i>	$D \leftarrow M[R[\%rsp]];$ $R[\%rsp] \leftarrow R[\%rsp] + 8$	Pop quad word

Figure 3.8 Push and pop instructions.

C declaration	Intel data type	Assembly-code suffix	Size (bytes)
char	Byte	b	1
short	Word	w	2
int	Double word	l	4
long	Quad word	q	8
char *	Quad word	q	8
float	Single precision	s	4
double	Double precision	l	8



Registers (X86_64)

- You get 16 registers

Type	Form	Operand value	Name
Immediate	$\$Imm$	Imm	Immediate
Register	r_a	$R[r_a]$	Register
Memory	Imm	$M[Imm]$	Absolute
Memory	(r_a)	$M[R[r_a]]$	Indirect
Memory	$Imm(r_b)$	$M[Imm + R[r_b]]$	Base + displacement
Memory	(r_b, r_i)	$M[R[r_b] + R[r_i]]$	Indexed
Memory	$Imm(r_b, r_i)$	$M[Imm + R[r_b] + R[r_i]]$	Indexed
Memory	$(, r_i, s)$	$M[R[r_i] \cdot s]$	Scaled indexed
Memory	$Imm(, r_i, s)$	$M[Imm + R[r_i] \cdot s]$	Scaled indexed
Memory	(r_b, r_i, s)	$M[R[r_b] + R[r_i] \cdot s]$	Scaled indexed
Memory	$Imm(r_b, r_i, s)$	$M[Imm + R[r_b] + R[r_i] \cdot s]$	Scaled indexed

Figure 3.3 Operand forms. Operands can denote immediate (constant) values, register values, or values from memory. The scaling factor s must be either 1, 2, 4, or 8.



ASM Maths

- inc/dec ++/--
- add/sub/imul
- xor/or/and
- AR shift
- Logical shift
- Variants for all sizes

Instruction		Effect	Description
leaq	S, D	$D \leftarrow \&S$	Load effective address
INC	D	$D \leftarrow D+1$	Increment
DEC	D	$D \leftarrow D-1$	Decrement
NEG	D	$D \leftarrow -D$	Negate
NOT	D	$D \leftarrow \sim D$	Complement
ADD	S, D	$D \leftarrow D + S$	Add
SUB	S, D	$D \leftarrow D - S$	Subtract
IMUL	S, D	$D \leftarrow D * S$	Multiply
XOR	S, D	$D \leftarrow D \wedge S$	Exclusive-or
OR	S, D	$D \leftarrow D \vee S$	Or
AND	S, D	$D \leftarrow D \& S$	And
SAL	k, D	$D \leftarrow D \ll k$	Left shift
SHL	k, D	$D \leftarrow D \ll k$	Left shift (same as SAL)
SAR	k, D	$D \leftarrow D \gg_A k$	Arithmetic right shift
SHR	k, D	$D \leftarrow D \gg_L k$	Logical right shift

ASM Maths: imult/div

- Requires more regs to avoid overflow
- Int128 types!

Instruction		Effect	Description
imulq	S	$R[\%rdx]:R[\%rax] \leftarrow S \times R[\%rax]$	Signed full multiply
mulq	S	$R[\%rdx]:R[\%rax] \leftarrow S \times R[\%rax]$	Unsigned full multiply
cqto		$R[\%rdx]:R[\%rax] \leftarrow \text{SignExtend}(R[\%rax])$	Convert to oct word
idivq	S	$R[\%rdx] \leftarrow R[\%rdx]:R[\%rax] \bmod S;$ $R[\%rax] \leftarrow R[\%rdx]:R[\%rax] \div S$	Signed divide
divq	S	$R[\%rdx] \leftarrow R[\%rdx]:R[\%rax] \bmod S;$ $R[\%rax] \leftarrow R[\%rdx]:R[\%rax] \div S$	Unsigned divide

Figure 3.12 Special arithmetic operations. These operations provide full 128-bit multiplication and division, for both signed and unsigned numbers. The pair of registers `%rdx` and `%rax` are viewed as forming a single 128-bit oct word.

ASM Control Flow

- Condition Codes

- Carry Flag

CF	(unsigned) t < (unsigned) a	Unsigned overflow
----	-----------------------------	-------------------

- Zero Flag

ZF	(t == 0)	Zero
----	----------	------

- Sign Flag

SF	(t < 0)	Negative
----	---------	----------

- Overflow

OF	(a < 0 == b < 0) && (t < 0 != a < 0)	Signed overflow
----	--------------------------------------	-----------------

ASM Control Flow

- CMP and TEST
- Results?

Instruction		Based on	Description
CMP	S_1, S_2	$S_2 - S_1$	Compare
cmpb			Compare byte
cmpw			Compare word
cmpl			Compare double word
cmpq			Compare quad word
TEST	S_1, S_2	$S_1 \& S_2$	Test
testb			Test byte
testw			Test word
testl			Test double word
testq			Test quad word

ASM Control Flow

- SET

Instruction	Synonym	Effect	Set condition
sete <i>D</i>	setz	$D \leftarrow ZF$	Equal / zero
setne <i>D</i>	setnz	$D \leftarrow \sim ZF$	Not equal / not zero
sets <i>D</i>		$D \leftarrow SF$	Negative
setns <i>D</i>		$D \leftarrow \sim SF$	Nonnegative
setg <i>D</i>	setnle	$D \leftarrow \sim (SF \wedge OF) \& \sim ZF$	Greater (signed >)
setge <i>D</i>	setnl	$D \leftarrow \sim (SF \wedge OF)$	Greater or equal (signed >=)
setl <i>D</i>	setnge	$D \leftarrow SF \wedge OF$	Less (signed <)
setle <i>D</i>	setng	$D \leftarrow (SF \wedge OF) \mid ZF$	Less or equal (signed <=)
seta <i>D</i>	setnbe	$D \leftarrow \sim CF \& \sim ZF$	Above (unsigned >)
setae <i>D</i>	setnb	$D \leftarrow \sim CF$	Above or equal (unsigned >=)
setb <i>D</i>	setnae	$D \leftarrow CF$	Below (unsigned <)
setbe <i>D</i>	setna	$D \leftarrow CF \mid ZF$	Below or equal (unsigned <=)

ASM Control Flow

- JMP
- C “goto”

<code>movq \$0,%rax</code>	<i>Set %rax to 0</i>
<code>jmp .L1</code>	<i>Goto .L1</i>
<code>movq (%rax),%rdx</code>	<i>Null pointer dereference (skipped)</i>
<code>.L1:</code>	
<code>popq %rdx</code>	<i>Jump target</i>

ASM Control Flow

- JMP operations

Instruction	Synonym	Jump condition	Description
jmp <i>Label</i>		1	Direct jump
jmp <i>*Operand</i>		1	Indirect jump
jz <i>Label</i>	jz	ZF	Equal / zero
jnz <i>Label</i>	jnz	~ZF	Not equal / not zero
js <i>Label</i>		SF	Negative
jns <i>Label</i>		~SF	Nonnegative
jg <i>Label</i>	jnle	~(SF ^ OF) & ~ZF	Greater (signed >)
jge <i>Label</i>	jnl	~(SF ^ OF)	Greater or equal (signed >=)
jl <i>Label</i>	jnge	SF ^ OF	Less (signed <)
jle <i>Label</i>	jng	(SF ^ OF) ZF	Less or equal (signed <=)
ja <i>Label</i>	jnbe	~CF & ~ZF	Above (unsigned >)
jae <i>Label</i>	jnb	~CF	Above or equal (unsigned >=)
jb <i>Label</i>	jnae	CF	Below (unsigned <)
jbe <i>Label</i>	jna	CF ZF	Below or equal (unsigned <=)

ASM Control Flow

- Putting it all together.

```
long lt_cnt = 0;
long ge_cnt = 0;

long absdiff_se(long x, long y)
{
    long result;
    if (x < y) {
        lt_cnt++;
        result = y - x;
    }
    else {
        ge_cnt++;
        result = x - y;
    }
    return result;
}

1 long gotodiff_se(long x, long y)
2 {
3     long result;
4     if (x >= y)
5         goto x_ge_y;
6     lt_cnt++;
7     result = y - x;
8     return result;
9 x_ge_y:
10    ge_cnt++;
11    result = x - y;
12    return result;
13 }
```

ASM Control Flow

- C Conversion

(b) Equivalent goto version

```
1 long gotodiff_se(long x, long y)
2 {
3     long result;
4     if (x >= y)
5         goto x_ge_y;
6     lt_cnt++;
7     result = y - x;
8     return result;
9 x_ge_y:
10    ge_cnt++;
11    result = x - y;
12    return result;
13 }
```

```
long absdiff_se(long x, long y)
x in %rdi, y in %rsi
1 absdiff_se:
2     cmpq    %rsi, %rdi           Compare x:y
3     jge     .L2                 If >= goto x_ge_y
4     addq    $1, lt_cnt(%rip)    lt_cnt++
5     movq    %rsi, %rax
6     subq    %rdi, %rax          result = y - x
7     ret                          Return
8 .L2:
9     addq    $1, ge_cnt(%rip)    x_ge_y:
10    movq    %rdi, %rax          ge_cnt++
11    subq    %rsi, %rax          result = x - y
12    ret                          Return
```

ASM Control Flow

- Branches and Moves

```
long absdiff(long x, long y)
{
    long result;
    if (x < y)
        result = y - x;
    else
        result = x - y;
    return result;
}
```

```
1  long cmovdiff(long x, long y)
2  {
3      long rval = y-x;
4      long eval = x-y;
5      long ntest = x >= y;
6      /* Line below requires
7         single instruction: */
8      if (ntest) rval = eval;
9      return rval;
10 }
```

ASM Control Flow

- C Conversion

long absdiff(long x, long y)

x in %rdi, y in %rsi

```
1 long cmovdiff(long x, long y)
2 {
3     long rval = y-x;
4     long eval = x-y;
5     long ntest = x >= y;
6     /* Line below requires
7        single instruction: */
8     if (ntest) rval = eval;
9     return rval;
10 }
```

```
1 absdiff:
```

```
2     movq    %rsi, %rax
```

```
3     subq    %rdi, %rax
```

rval = y-x

```
4     movq    %rdi, %rdx
```

```
5     subq    %rsi, %rdx
```

eval = x-y

```
6     cmpq    %rsi, %rdi
```

Compare x:y

```
7     cmovge  %rdx, %rax
```

If >=, rval = eval

```
8     ret
```

Return tval

ASM Control Flow

- CMOV instructions

Instruction		Synonym	Move condition	Description
<code>cmovz</code>	S, R	<code>cmovz</code>	ZF	Equal / zero
<code>cmovne</code>	S, R	<code>cmovnz</code>	$\sim$ ZF	Not equal / not zero
<code>cmovs</code>	S, R		SF	Negative
<code>cmovns</code>	S, R		$\sim$ SF	Nonnegative
<code>cmovg</code>	S, R	<code>cmovnle</code>	$\sim(SF \wedge OF) \ \& \ \sim ZF$	Greater (signed $>$)
<code>cmovge</code>	S, R	<code>cmovnl</code>	$\sim(SF \wedge OF)$	Greater or equal (signed $\geq$)
<code>cmovl</code>	S, R	<code>cmovnge</code>	$SF \wedge OF$	Less (signed $<$)
<code>cmovle</code>	S, R	<code>cmovng</code>	$(SF \wedge OF) \mid ZF$	Less or equal (signed $\leq$)
<code>cmova</code>	S, R	<code>cmovnbe</code>	$\sim CF \ \& \ \sim ZF$	Above (unsigned $>$)
<code>cmovae</code>	S, R	<code>cmovnb</code>	$\sim CF$	Above or equal (Unsigned $\geq$)
<code>cmovb</code>	S, R	<code>cmovnae</code>	CF	Below (unsigned $<$)
<code>cmovbe</code>	S, R	<code>cmovna</code>	CF $\mid$ ZF	Below or equal (unsigned $\leq$)

Figure 3.18 The conditional move instructions. These instructions copy the source value S to its destination R when the move condition holds. Some instructions have “synonyms,” alternate names for the same machine instruction.

ASM Control Flow

- Loops

```
long fact_do(long n)
{
    long result = 1;
    do {
        result *= n;
        n = n-1;
    } while (n > 1);
    return result;
}
```

```
long fact_do_goto(long n)
{
    long result = 1;
loop:
    result *= n;
    n = n-1;
    if (n > 1)
        goto loop;
    return result;
}
```

ASM Control Flow

- C Conversion

```
long fact_do_goto(long n)
{
    long result = 1;
loop:
    result *= n;
    n = n-1;
    if (n > 1)
        goto loop;
    return result;
}
```

```
long fact_do(long n)
n in %rdi
1  fact_do:
2      movl    $1, %eax           Set result = 1
3      .L2:                       loop:
4      imulq   %rdi, %rax         Compute result *= n
5      subq    $1, %rdi           Decrement n
6      cmpq    $1, %rdi           Compare n:1
7      jg      .L2                If >, goto loop
8      rep; ret                  Return
```

ASM Control Flow

- Loops

```
long fact_while(long n)
{
    long result = 1;
    while (n > 1) {
        result *= n;
        n = n-1;
    }
    return result;
}
```

```
long fact_while_jm_goto(long n)
{
    long result = 1;
    goto test;
loop:
    result *= n;
    n = n-1;
test:
    if (n > 1)
        goto loop;
    return result;
}
```

ASM Control Flow

- C Conversion

```
long fact_while_jm_goto(long n)
{
    long result = 1;
    goto test;
loop:
    result *= n;
    n = n-1;
test:
    if (n > 1)
        goto loop;
    return result;
}
```

```
long fact_while(long n)
n in %rdi
fact_while:
    movl    $1, %eax        Set result = 1
    jmp     .L5             Goto test
.L6:
    imulq   %rdi, %rax       Compute result *= n
    subq    $1, %rdi        Decrement n
.L5:
    cmpq    $1, %rdi        Compare n:1
    jg      .L6             If >, goto loop
    rep; ret                Return
```

ASM Control Flow

- Case/Switch

```
void switch_eg(long x, long n,
               long *dest)
{
    long val = x;

    switch (n) {

        case 100:
            val *= 13;
            break;

        case 102:
            val += 10;
            /* Fall through */

        case 103:
            val += 11;
            break;

        case 104:
        case 106:
            val *= val;
            break;

        default:
            val = 0;
    }
    *dest = val;
}
```

```
1 void switch_eg_impl(long x, long n,
2                     long *dest)
3 {
4     /* Table of code pointers */
5     static void *jt[7] = {
6         &loc_A, &loc_def, &loc_B,
7         &loc_C, &loc_D, &loc_def,
8         &loc_D
9     };
10    unsigned long index = n - 100;
11    long val;
12
13    if (index > 6)
14        goto loc_def;
15    /* Multiway branch */
16    goto *jt[index];
17
18    loc_A:    /* Case 100 */
19        val = x * 13;
20        goto done;
21    loc_B:    /* Case 102 */
22        x = x + 10;
23        /* Fall through */
24    loc_C:    /* Case 103 */
25        val = x + 11;
26        goto done;
27    loc_D:    /* Cases 104, 106 */
28        val = x * x;
29        goto done;
30    loc_def:  /* Default case */
31        val = 0;
32    done:
33        *dest = val;
34 }
```

ASM Control Flow

- C Conversion

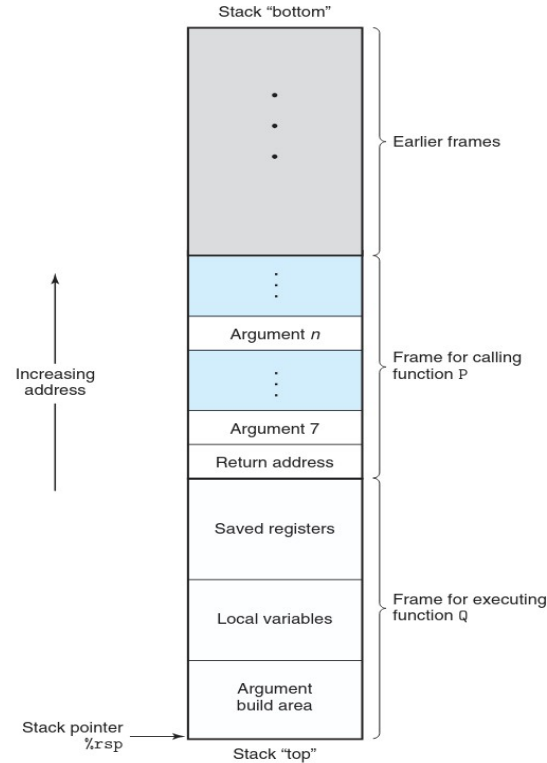
```
void switch_eg(long x, long n, long *dest)
x in %rdi, n in %rsi, dest in %rdx

1  switch_eg:
2      subq    $100, %rsi           Compute index = n-100
3      cmpq    $6, %rsi            Compare index:6
4      ja      .L8                 If >, goto loc_def
5      jmp     *.L4(,%rsi,8)        Goto *jg[index]
6  .L3:
7      leaq    (%rdi,%rdi,2), %rax   3*x
8      leaq    (%rdi,%rax,4), %rdi   val = 13*x
9      jmp     .L2                 Goto done
10 .L5:
11      addq    $10, %rdi           x = x + 10
12 .L6:
13      addq    $11, %rdi           val = x + 11
14      jmp     .L2                 Goto done
15 .L7:
16      imulq   %rdi, %rdi          val = x * x
17      jmp     .L2                 Goto done
18 .L8:
19      movl    $0, %edi            val = 0
20 .L2:
21      movq    %rdi, (%rdx)        *dest = val
22      ret                        Return
```

ASM Procedures

- Functions?
- More STACKS

Instruction	Description
<code>call <i>Label</i></code>	Procedure call
<code>call *<i>Operand</i></code>	Procedure call
<code>ret</code>	Return from call



ASM Procedures

- Functions?

```
void proc(long a1, long *a1p,  
          int a2, int *a2p,  
          short a3, short *a3p,  
          char a4, char *a4p)  
{  
    *a1p += a1;  
    *a2p += a2;  
    *a3p += a3;  
    *a4p += a4;  
}
```

```
1  proc:  
2      movq    16(%rsp), %rax    Fetch a4p (64 bits)  
3      addq    %rdi, (%rsi)     *a1p += a1 (64 bits)  
4      addl    %edx, (%rcx)     *a2p += a2 (32 bits)  
5      addw    %r8w, (%r9)      *a3p += a3 (16 bits)  
6      movl    8(%rsp), %edx     Fetch a4 ( 8 bits)  
7      addb    %dl, (%rax)      *a4p += a4 ( 8 bits)  
8      ret                               Return
```

ASM Procedures

- Calling a function

```
long call_proc()
{
    long x1 = 1; int x2 = 2;
    short x3 = 3; char x4 = 4;
    proc(x1, &x1, x2, &x2, x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

```
long call_proc()
1  call_proc:
    Set up arguments to proc
2      subq    $32, %rsp        Allocate 32-byte stack frame
3      movq    $1, 24(%rsp)     Store 1 in &x1
4      movl    $2, 20(%rsp)     Store 2 in &x2
5      movw    $3, 18(%rsp)     Store 3 in &x3
6      movb    $4, 17(%rsp)     Store 4 in &x4
7      leaq    17(%rsp), %rax    Create &x4
8      movq    %rax, 8(%rsp)     Store &x4 as argument 8
9      movl    $4, (%rsp)       Store 4 as argument 7
10     leaq    18(%rsp), %r9     Pass &x3 as argument 6
11     movl    $3, %r8d          Pass 3 as argument 5
12     leaq    20(%rsp), %rcx    Pass &x2 as argument 4
13     movl    $2, %edx          Pass 2 as argument 3
14     leaq    24(%rsp), %rsi    Pass &x1 as argument 2
15     movl    $1, %edi          Pass 1 as argument 1
    Call proc
16     call    proc
    Retrieve changes to memory
17     movslq   20(%rsp), %rdx    Get x2 and convert to long
18     addq     24(%rsp), %rdx    Compute x1+x2
19     movswl   18(%rsp), %eax    Get x3 and convert to int
20     movsbl   17(%rsp), %ecx    Get x4 and convert to int
21     subl     %ecx, %eax        Compute x3-x4
22     cltq                                           Convert to long
23     imulq    %rdx, %rax        Compute (x1+x2) * (x3-x4)
24     addq     $32, %rsp        Deallocate stack frame
25     ret                                           Return
```

ASM Procedures

- Recursive Calls

```
long rfact(long n)
{
    long result;
    if (n <= 1)
        result = 1;
    else
        result = n * rfact(n-1);
    return result;
}
```

```
long rfact(long n)
n in %rdi
1  rfact:
2      pushq    %rbx                Save %rbx
3      movq     %rdi, %rbx         Store n in callee-saved register
4      movl     $1, %eax           Set return value = 1
5      cmpq     $1, %rdi           Compare n:1
6      jle      .L35               If <=, goto done
7      leaq     -1(%rdi), %rdi     Compute n-1
8      call     rfact              Call rfact(n-1)
9      imulq    %rbx, %rax         Multiply result by n
10     .L35:                       done:
11     popq     %rbx               Restore %rbx
12     ret                                Return
```

ASM Procedures

- Arrays(Size of LONG?)

```
long P[M][N];
```

```
long Q[N][M];
```

```
long sum_element(long i, long j) {  
    return P[i][j] + Q[j][i];  
}
```

```
long sum_element(long i, long j)
```

```
i in %rdi, j in %rsi
```

```
1  sum_element:  
2      leaq    0(,%rdi,8), %rdx  
3      subq    %rdi, %rdx  
4      addq    %rsi, %rdx  
5      leaq    (%rsi,%rsi,4), %rax  
6      addq    %rax, %rdi  
7      movq    Q(,%rdi,8), %rax  
8      addq    P(,%rdx,8), %rax  
9      ret
```

ASM Floats

- Even more regs

Instruction	Source	Destination	Description
<code>vmovss</code>	M_{32}	X	Move single precision
<code>vmovss</code>	X	M_{32}	Move single precision
<code>vmovsd</code>	M_{64}	X	Move double precision
<code>vmovsd</code>	X	M_{64}	Move double precision
<code>vmovaps</code>	X	X	Move aligned, packed single precision
<code>vmovapd</code>	X	X	Move aligned, packed double precision

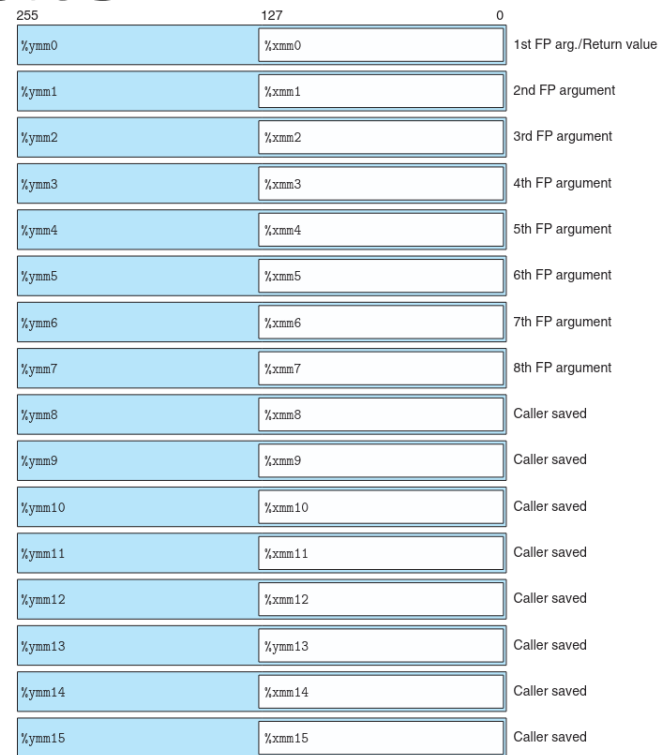


Figure 3.45 Media registers. These registers are used to hold floating-point data. Each YMM register holds 32 bytes. The low-order 16 bytes can be accessed as an XMM register.

ASM Floats

- Conversions

Instruction	Source	Destination	Description
<code>vcvttss2si</code>	X/M_{32}	R_{32}	Convert with truncation single precision to integer
<code>vcvttsd2si</code>	X/M_{64}	R_{32}	Convert with truncation double precision to integer
<code>vcvttss2siq</code>	X/M_{32}	R_{64}	Convert with truncation single precision to quad word integer
<code>vcvttsd2siq</code>	X/M_{64}	R_{64}	Convert with truncation double precision to quad word integer

Figure 3.47 Two-operand floating-point conversion operations. These convert floating-point data to integers. (X : XMM register (e.g., `%xmm3`); R_{32} : 32-bit general-purpose register (e.g., `%eax`); R_{64} : 64-bit general-purpose register (e.g., `%rax`); M_{32} : 32-bit memory range; M_{64} : 64-bit memory range)

ASM Floats

- Conversions

Instruction	Source 1	Source 2	Destination	Description
<code>vcvtsi2ss</code>	M_{32}/R_{32}	X	X	Convert integer to single precision
<code>vcvtsi2sd</code>	M_{32}/R_{32}	X	X	Convert integer to double precision
<code>vcvtsi2ssq</code>	M_{64}/R_{64}	X	X	Convert quad word integer to single precision
<code>vcvtsi2sdq</code>	M_{64}/R_{64}	X	X	Convert quad word integer to double precision

Figure 3.48 Three-operand floating-point conversion operations. These instructions convert from the data type of the first source to the data type of the destination. The second source value has no effect on the low-order bytes of the result. (X : XMM register (e.g., `%xmm3`); M_{32} : 32-bit memory range; M_{64} : 64-bit memory range)

ASM Floats

- C to ASM

```
double fcvt(int i, float *fp, double *dp, long *lp)
{
    float f = *fp; double d = *dp; long l = *lp;
    *lp = (long)    d;
    *fp = (float)   i;
    *dp = (double)  l;
    return (double) f;
}
```

```
double fcvt(int i, float *fp, double *dp, long *lp)
i in %edi, fp in %rsi, dp in %rdx, lp in %rcx
1  fcvt:
2      vmovss  (%rsi), %xmm0           Get f = *fp
3      movq    (%rcx), %rax           Get l = *lp
4      vcvttss2siq  (%rdx), %r8       Get d = *dp and convert to long
5      movq    %r8, (%rcx)           Store at lp
6      vcvtsi2ss    %edi, %xmm1, %xmm1  Convert i to float
7      vmovss  %xmm1, (%rsi)         Store at fp
8      vcvtsi2sdq    %rax, %xmm1, %xmm1  Convert l to double
9      vmovsd  %xmm1, (%rdx)         Store at dp
    The following two instructions convert f to double
10     vunpcklps    %xmm0, %xmm0, %xmm0
11     vcvtps2pd    %xmm0, %xmm0
12     ret                               Return f
```

ASM Floats

- Even more math

Single	Double	Effect	Description
vaddss	vaddsd	$D \leftarrow S_2 + S_1$	Floating-point add
vsubss	vsubsd	$D \leftarrow S_2 - S_1$	Floating-point subtract
vmulss	vmulsd	$D \leftarrow S_2 \times S_1$	Floating-point multiply
vdivss	vdivsd	$D \leftarrow S_2 / S_1$	Floating-point divide
vmaxss	vmaxsd	$D \leftarrow \max(S_2, S_1)$	Floating-point maximum
vminss	vminsd	$D \leftarrow \min(S_2, S_1)$	Floating-point minimum
sqrtps	sqrtsd	$D \leftarrow \sqrt{S_1}$	Floating-point square root

ASM Floats

- Comparison and Control

Instruction		Based on	Description
ucomiss	S_1, S_2	$S_2 - S_1$	Compare single precision
ucomisd	S_1, S_2	$S_2 - S_1$	Compare double precision

Ordering $S_2:S_1$	CF	ZF	PF
Unordered	1	1	1
$S_2 < S_1$	1	0	0
$S_2 = S_1$	0	1	0
$S_2 > S_1$	0	0	0

ASM Floats

- C to ASM Control

```
typedef enum {NEG, ZERO, POS, OTHER} range_t;
```

```
range_t find_range(float x)
{
    int result;
    if (x < 0)
        result = NEG;
    else if (x == 0)
        result = ZERO;
    else if (x > 0)
        result = POS;
    else
        result = OTHER;
    return result;
}
```

```
range_t find_range(float x)
x in %xmm0

1 find_range:
2   vxorps  %xmm1, %xmm1, %xmm1
3   vucomiss %xmm0, %xmm1
4   ja     .L5
5   vucomiss %xmm1, %xmm0
6   jp     .L8
7   movl    $1, %eax
8   je     .L3
9   .L8:
10  vucomiss .LC0(%rip), %xmm0
11  setbe    %al
12  movzbl   %al, %eax
13  addl     $2, %eax
14  ret
15  .L5:
16  movl     $0, %eax
17  .L3:
18  rep; ret

Set %xmm1 = 0
Compare 0:x
If >, goto neg
Compare x:0
If NaN, goto posornan
result = ZERO
If =, goto done
posornan:
Compare x:0
Set result = NaN ? 1 : 0
Zero-extend
result += 2 (POS for > 0, OTHER for NaN)
Return
neg:
result = NEG
done:
Return
```